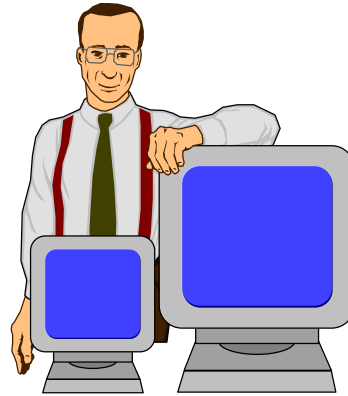


ATM Traffic Management

Dollar Day Sale



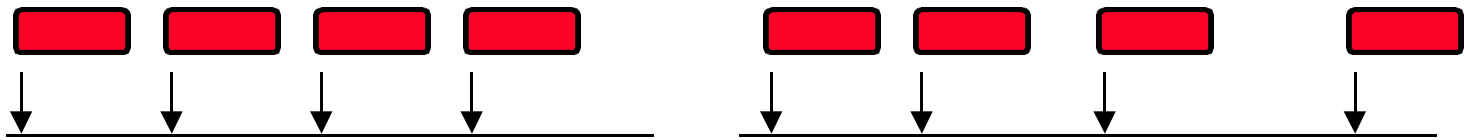
One Megabit memory, One Megabyte disk,
One Mbps link, One MIP processor, one
dollar each.....

Raj Jain

**Raj Jain is now at
Washington University in Saint Louis
Jain@cse.wustl.edu
<http://www.cse.wustl.edu/~jain/>**

ATM Networks: Overview

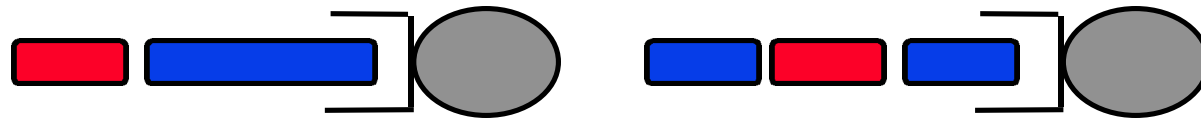
- q STM = Synchronous Transfer Mode,
ATM = Asynchronous Transfer Mode



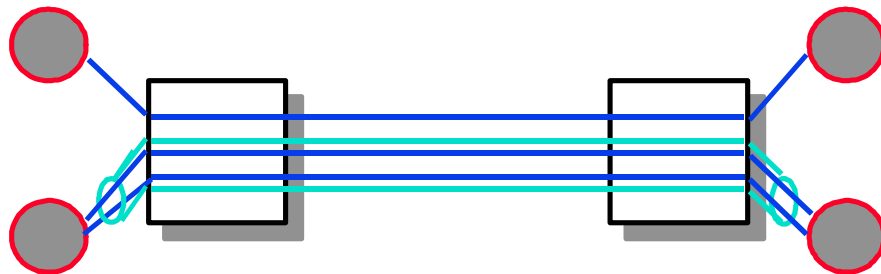
Allows **any-speed** and even **variable rate** connection

Broadband = Rate greater than primary rate (1.5 Mbps)

- q ATM = Short fixed size 53-byte cells



- q Connection oriented \Rightarrow Virtual Channels (VC)





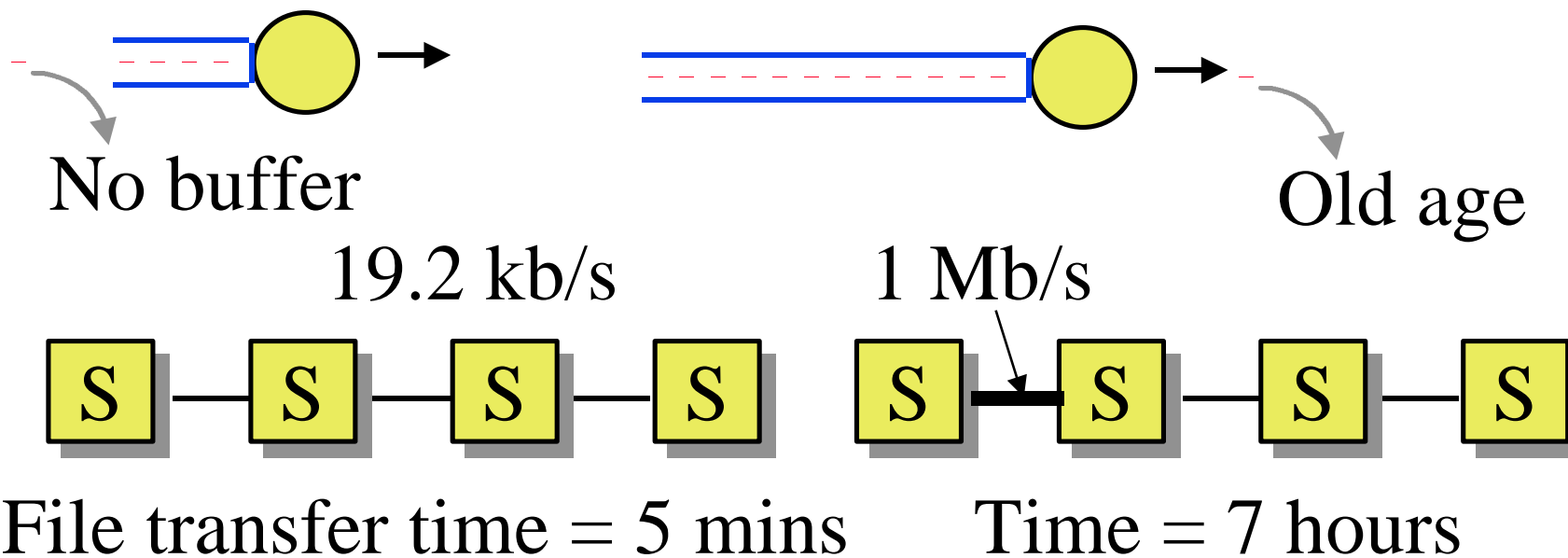
- q Why worry about congestion?
- q Congestion schemes for ATM
- q Rate vs Credit: Key issues
- q Explicit Rate-based Control
- q ABR Traffic Management

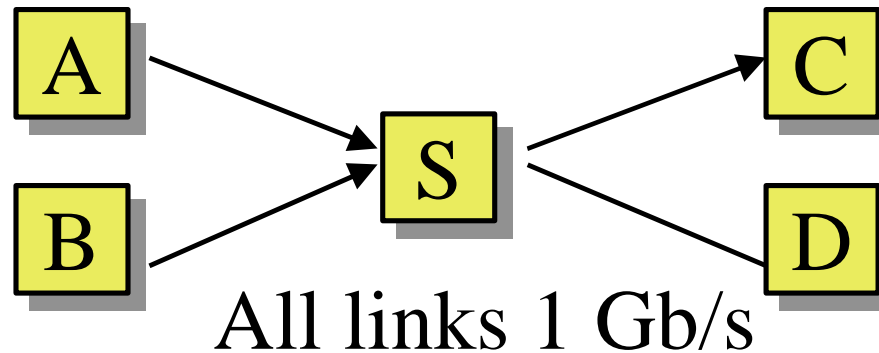
Why Worry About Congestion?

Q: Will the congestion problem be solved when:

- q Memory becomes cheap (infinite memory)?
- q Links become cheap (very high speed links)?
- q Processors become cheap?

A: None of the above.





Conclusions:

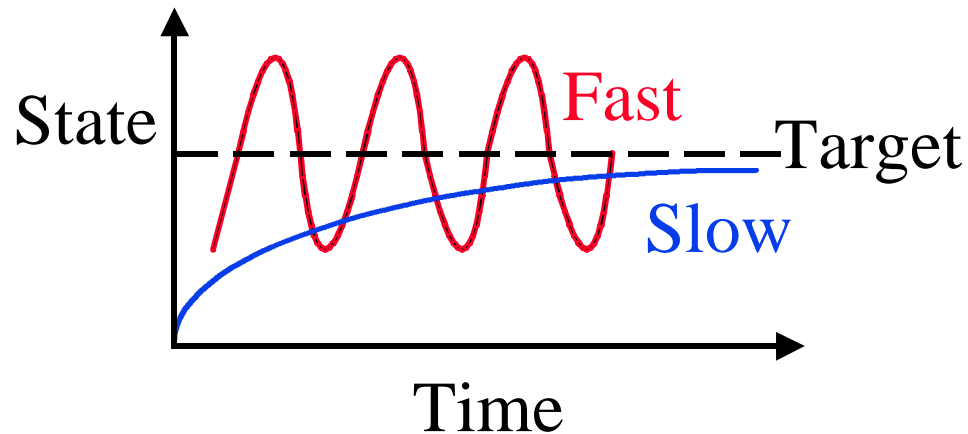
- q Congestion is a dynamic problem.
Static solutions are not sufficient
- q Bandwidth explosion
⇒ More unbalanced networks
- q Buffer shortage is a symptom not the cause.

Economic Reasons

- q Network is a shared resource
Because it is expensive and needed occasionally
(Like airplanes, emergency rooms)
- q Most costs are fixed.
Cost for fiber, switches, laying fiber and
maintaining them does not depend upon usage
⇒ Underutilization is expensive
- q But overutilization leads to user dissatisfaction.
- q Need a way to keep the network maximally utilized

One Scheme or Many?

- q Fundamental principle of control theory:



- q Control faster than feedback \Rightarrow Instability
- Control slower than feedback \Rightarrow non-responsiveness
- Ideal: Control rate \approx Feedback rate
- q Lesson: No scheme can cure a congestion lasting less than its feedback delay.

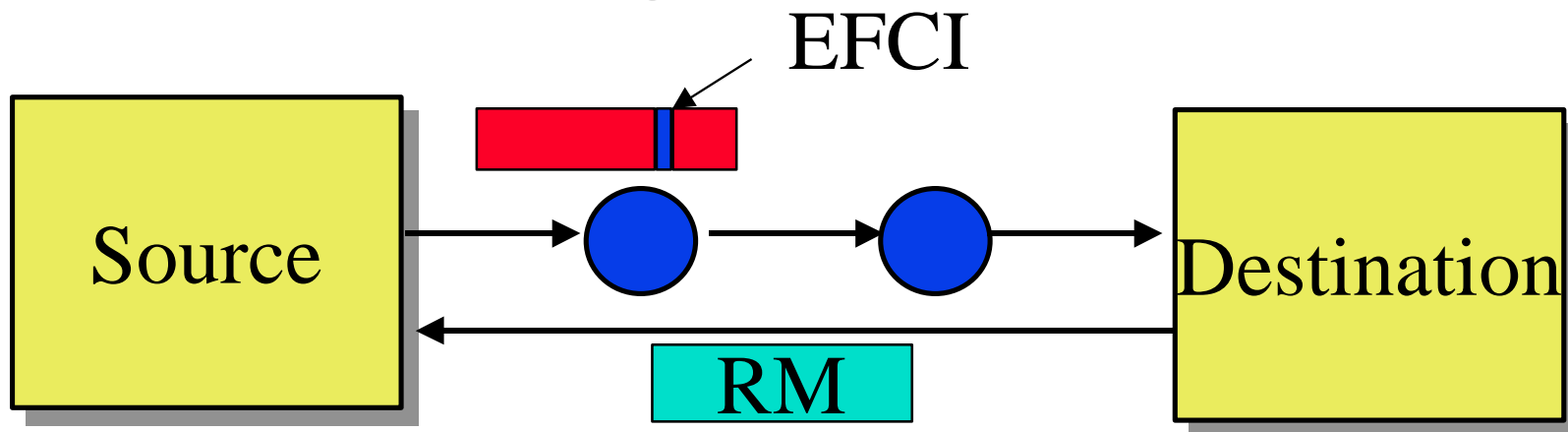
Classes of Service

- q **CBR** (Constant bit rate): User declares required rate. Throughput, delay and delay variation guaranteed.
- q **VBR** (Variable bit rate): User declares average and max rate.
 - **rt-VBR** (Real-time variable bit rate): Conferencing. Max delay and delay variation guaranteed.
 - **nrt-VBR** (non-real time variable bit rate): Stored video.
- q **ABR** (Available bit rate): Follows feedback instructions. Network gives maximum throughput with minimum loss.
- q **UBR** (Unspecified bit rate): User sends whenever it wants. No feedback mechanism. No guarantee. Cells may be dropped during congestion.

Traffic Management Functions

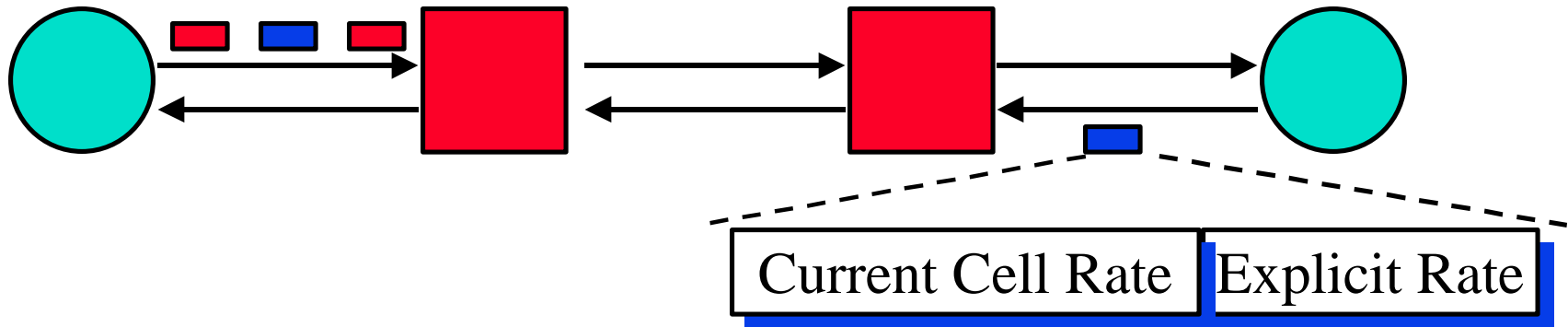
- q Connection Admission Control (CAC):
Can requested bandwidth and quality of service be supported?
- q Traffic Shaping: Limit burst length. Space-out cells.
- q Usage Parameter Control (UPC):
Monitor and control traffic at the network entrance.
- q Network Resource Management:
Scheduling, Queueing, virtual path resource reservation
- q Selective cell discard:
Cell Loss Priority (CLP) = 1 cells may be dropped
Cells of non-compliant connections may be dropped
- q Frame Discarding
- q Feedback Controls: Network tells the source to increase or decrease its load.

Initial Binary Rate-based Scheme



- q *We invented DECbit scheme (1986).
Implemented in many standards since 1986.*
- q *Forward explicit congestion notification (FECN) in Frame relay*
- q *Explicit forward congestion indicator (EFCI) set to 0 at source.
Congested switches set EFCI to 1*
- q *Every n th cell, destination sends an resource management (RM) cell to the source indicating increase amount or decrease factor*

The Explicit Rate Scheme

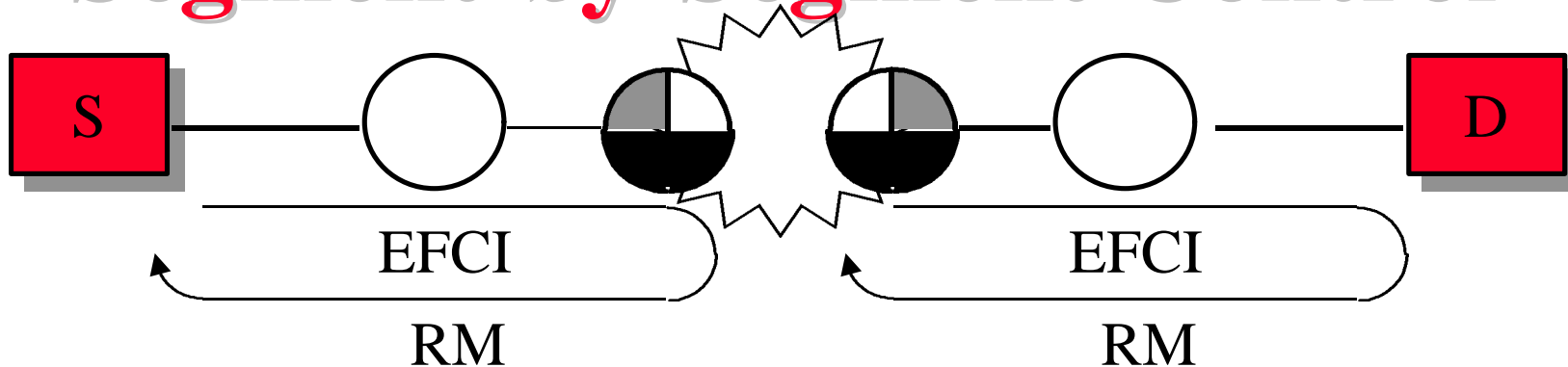


- q Sources send one **RM cell** every n cells
- q The RM cells contain “**Explicit rate**”
- q Destination returns the RM cell to the source
- q The switches adjust the rate **down**
- q Source adjusts to the specified rate

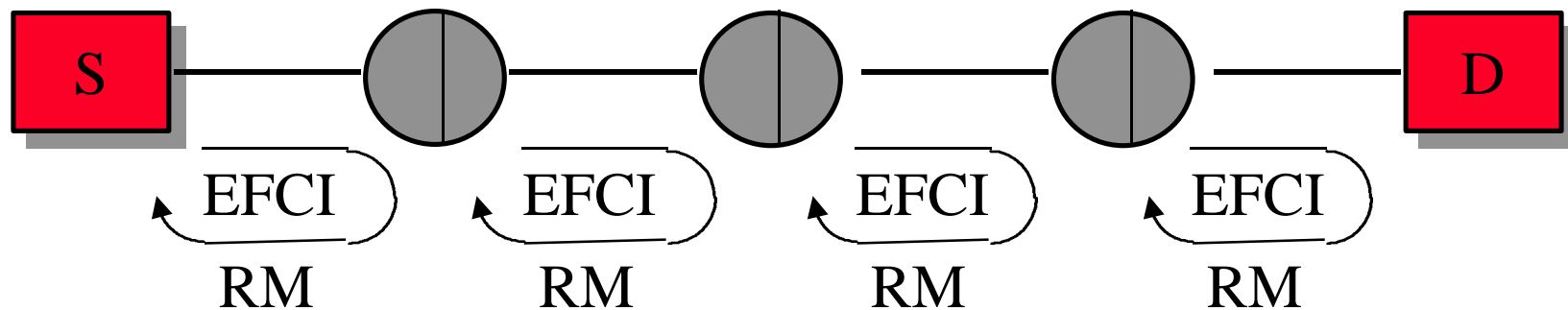
ERICA Switch Algorithm

- q Each manufacturer will have its own explicit rate switch algorithm
- q Explicit Rate Indication for Congestion Avoidance (ERICA) is the most thoroughly analyzed algorithm among disclosed algorithms
- q Shown to be efficient, fair, fast transient response, able to handle bursty TCP traffic
- q ERICA+ allows low delay even at 100% utilization and provides stability in the presence of high frequency VBR background traffic
- q Being implemented by several vendors

Segment-by-Segment Control

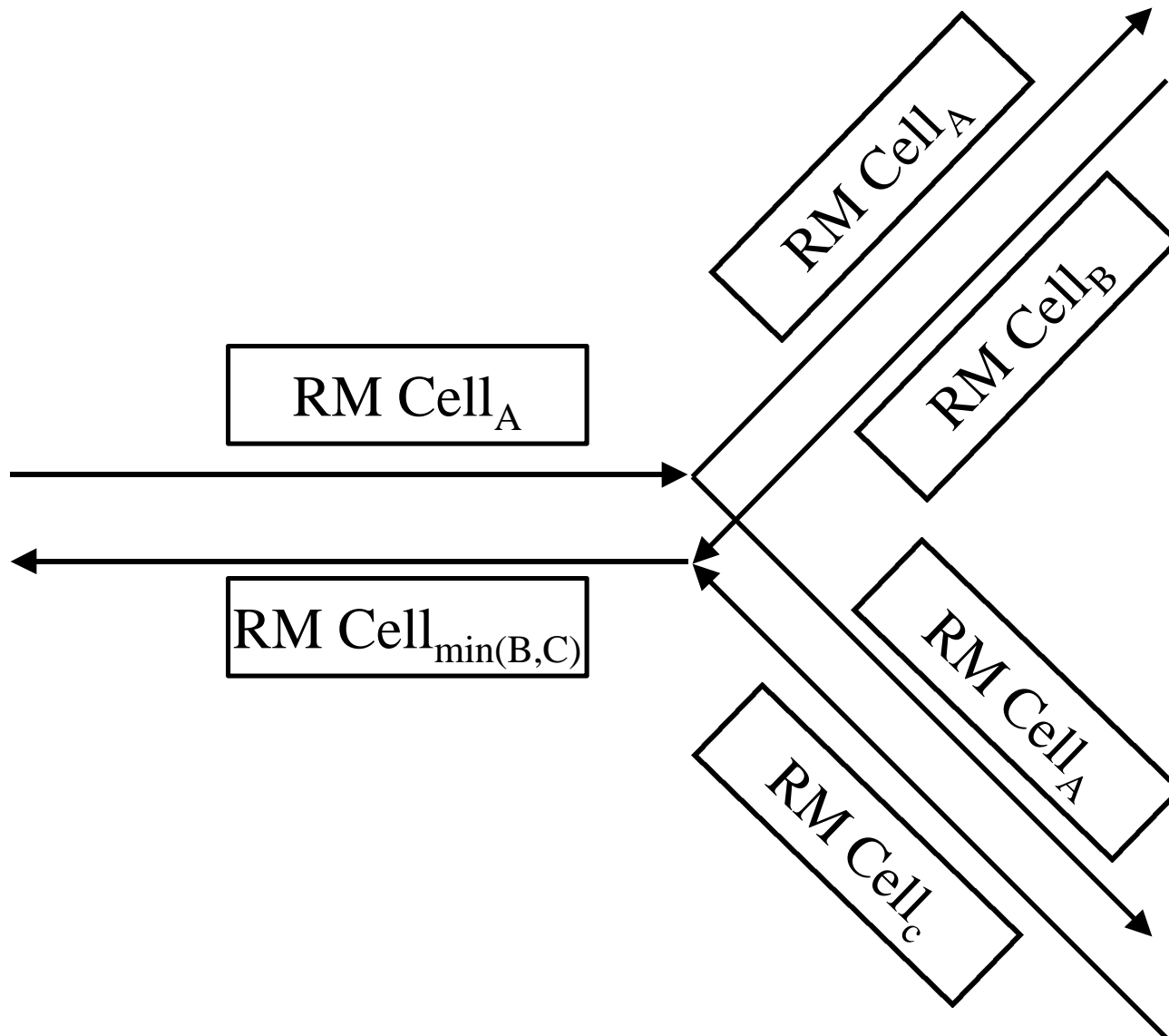


- q Virtual source/virtual destinations follow all notification/control rules
- q Can be hop-by-hop



- q Virtual dest/sources maintain per-VC queues.

Multicast

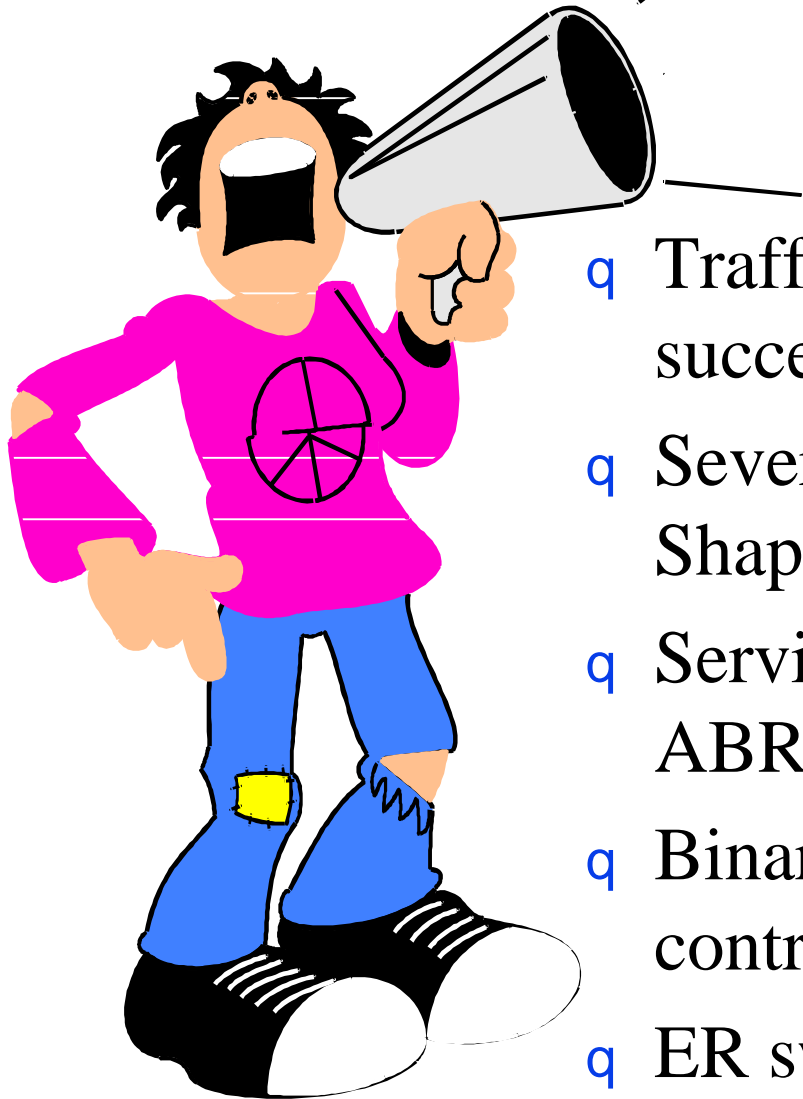


Outstanding Issues

- q Bursty sources: Client server, transactions, WWW
- q Effect of parameters: Optimal parameter values
- q Priority service for RM cells
- q Multicast
- q Connection admission control (CAC)
- q TCP/IP over UBR
- q Non-conforming sources
- q Optimal Source Strategy: Parameter + Out-of-rate cells
- q Virtual Source/destination
- q Implicit feedback schemes: Heterogeneous Networks



Congestion: Summary



- q Traffic Management is key to success of ATM
- q Several different methods: CAC, Shaping, UPC, Scheduling, ...
- q Service categories: CBR, VBR, ABR, UBR
- q Binary feedback too slow for rate control. Especially for satellites.
- q ER switches provide much better performance than EFCI.

References

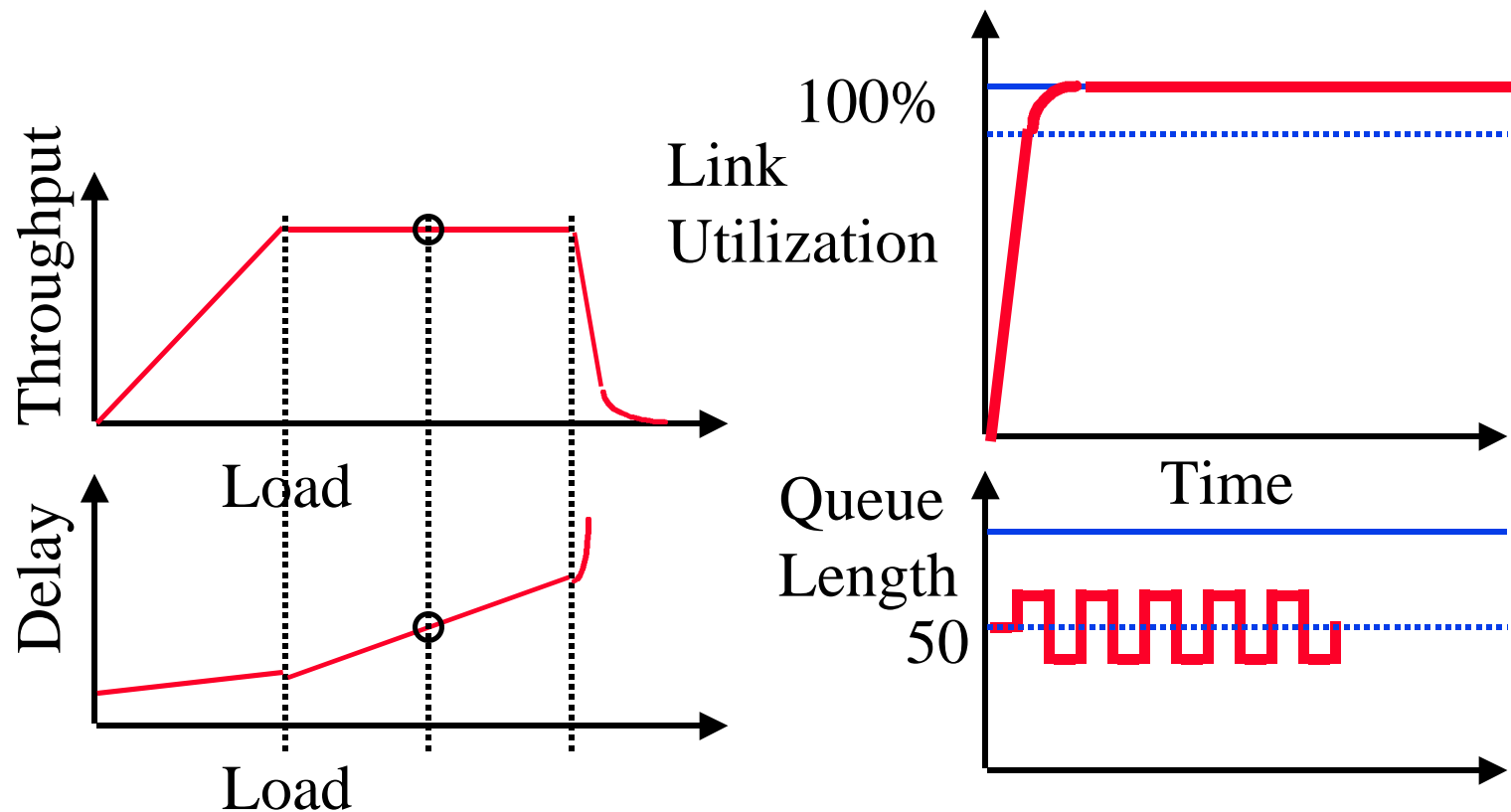
- q R.Jain, "Congestion Control and Traffic Management in ATM Networks: Recent Advances and A Survey", Invited submission to Computer Networks and ISDN Systems, February 1995.
Available at <http://www.cis.ohio-state.edu/~jain/>
- q *User-Network Interface Specifications, V3.0*, Prentice-Hall, September 10, 1993., (515)-284-6751
V3.1 released by ATM Forum, January 1995.
- q K. Siu and R. Jain, "A Brief Overview of ATM: Protocol Layers, LAN Emulation, and Traffic Management," Computer Communications Review (ACM SIGCOMM), April 1995.
- q "ATM Forum Traffic Management Specification, Version 4.0," ATM Forum/95-0013R13, April 1996.

- q Craig Partridge, *Gigabit Networking*, Addison-Wesley, 1993.
- q S. Akhtar, *Congestion Control in a Fast Packet Switching Network*, Master's Thesis, Washington University, 1987.
[Originated the concept of leaky bucket.]
- q Allyn Romanow, *Packet Discard Strategies for Controlling AM Congestion*, ATM Forum Traffic Management Subworking Group submission #94-0107, January 17-24, 1994.
- q S. Kalyanaraman, R. Jain, S. Fahmy, R. Goyal, F. Lu and S. Srinidhi, "Performance of TCP/IP over ABR," Submitted to Globecom'96.
- q R. Jain, S. Kalyanaraman, R. Viswanathan, "The OSU Scheme for Congestion Avoidance in ATM networks Using Explicit Rate Indication," Proceedings WATM'95, Paris, December 1995, Also available an extended OSU technical report and an extended conference version.

- q "Killer App? It's the high-speed backbone," An interview with Raj Jain. Government Computer News, October 16, 1995, pp. 24.
- q R. Jain, "ABR Service on ATM Networks: What is it?" Network World, June 24, 1995.
- q A. Charny, D. Clark, R. Jain, "Congestion Control with Explicit Rate Indication," Proc. ICC'95, June 1995, 10 pp.

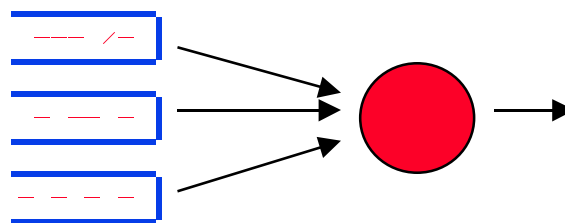
ERICA+: Full Utilization

- q Allows operation at any point between the knee and the cliff
- q The queue time can be set to any desired value.
- q Allows utilization to be 100%



Fair Queueing

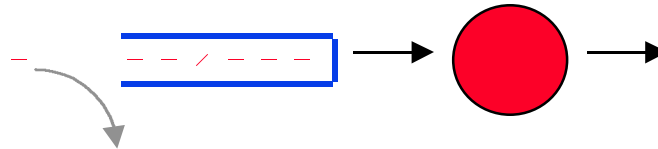
- q Goal: Isolation of flows.
- q Round Robin:



- q Problem: Flows with large packets get more bandwidth.
- q Bit-by-bit Round Robin
- q Fair Queueing: Compute end-times of packets under bit-by-bit round robin.
- q Weighted Fair Queueing: Different flows are allowed different bandwidth.

Selective Cell Discarding

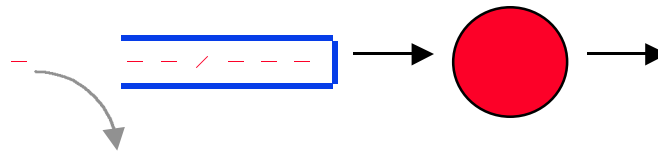
- q **Goal:** Discard cells of VCs not following contract
- q **Normal Discard:** CLP=1 cells are discarded if buffer is full.



- q **Selective Cell Discard:** May discard CLP =1 or CLP =0 cells of non-compliant VCs.

Early Packet Discard

- q **Goal:** To minimize packet loss rate
- q **Normal Discard:** Cells are discarded if buffer is full.



Problem: Unfair to late-comers.

- q **Random Discard:** Randomly select a cell in the queue and discard it.

Problem: Need to look inside the queues.

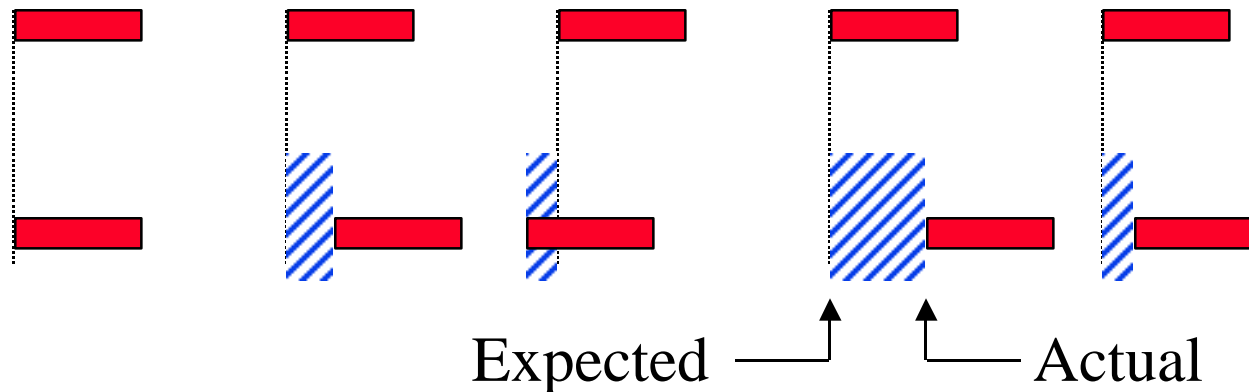
- q **Early Packet Discard:** When a queue threshold exceeds,
 - First cell of the next arriving packet is discarded.
 - Drop all cells of discarded packets.

Cell Delay Variation

- q Cell Transfer Delay (CTD): First bit out to last bit in
- q Cell Delay Variation (CDV) = $CTD_{\max} - CTD_{\min}$
 - Peak-to-peak CDV
 - Instantaneous CDV

Instantaneous CDV

- q I-CDV = Actual - Expected arrival time
- q Expected = Emission + Nominal delay
- q Cell Delay Variation Window (CDV-W)
 $CDV-W = |I-CDV(Max)| + |I-CDV(Min)|$
- q Cells arriving outside window are considered lost
- q Large CDV \Rightarrow Large buffers \Rightarrow Higher cost



Generic Flow Control (GFC)

- q ATM header format designed from DQDB cell headers. The first byte in DQDB is access control for contention.
- q Some countries wanted to keep this in ATM for compatibility and others saw no use for it.
- q Compromise \Rightarrow One half of the octet was reserved for “Generic Flow Control” to be defined at some future point.
- q British Telecom adapted multiservice flow control (MSFC) from Orwell ring and proposed to use it for GFC

Fast Resource Management

- q User requests to send a burst.
- q Network reserves the resources and grants the request
- q User sends the bursts
- q If network rejects, ???

ERICA Features

- q Measured overload/load at switch
- q Small queue lengths during steady state
- q Fast response due to optimistic design
- q Parameters: Few, insensitive, easy
- q Insensitive to source not using their allocated rates
- q Several options: Backward Explicit Congestion Notification
- q Simplified switch algorithm
- q Optimized all steps. Eliminated unnecessary steps. Eliminated many parameters

ERICA+: Switch Algorithm

- q Target cell rate = Target Utilization \times Link Capacity
- q Target Utilization
= fn(Current load, Queue length, Queue drain time goal)
- q Rest is similar to ERICA
- q Features:
 - Queue length is bounded during overload
 - No queue underflow \Rightarrow Switches keep ABR cells waiting to be transmitted as soon as the bandwidth becomes available.
 - 100% Utilization even with VBR

Control Mechanisms

- q nrt-VBR: Open-Loop + Optional closed loop component
 - Traffic shaping, CLP
 - EFCI optional
 - ⇒ No switch or end-system behavior specified
- q UBR: Local policy. CLP=0 or 1. Not subject to CAC.
- q ABR:
 - CLP=1 data cells not allowed.
 - No EFCI in RM cells
 - All parameters are negotiated independently for the two directions
 - MCR=0 not subject to CAC due to load

Rate Representation

Reserved	Nonzero	Exponent	Mantissa
----------	---------	----------	----------

1 1 5 9 ← Size in bits

q Rate in cells/second = $[2^e(1+m/512)] * nz$

q Example:

$$0-1-01010-0\ 1100\ 1010 = 2^{18}(1+202/512)$$

$$= 262144 \times 1.394523$$

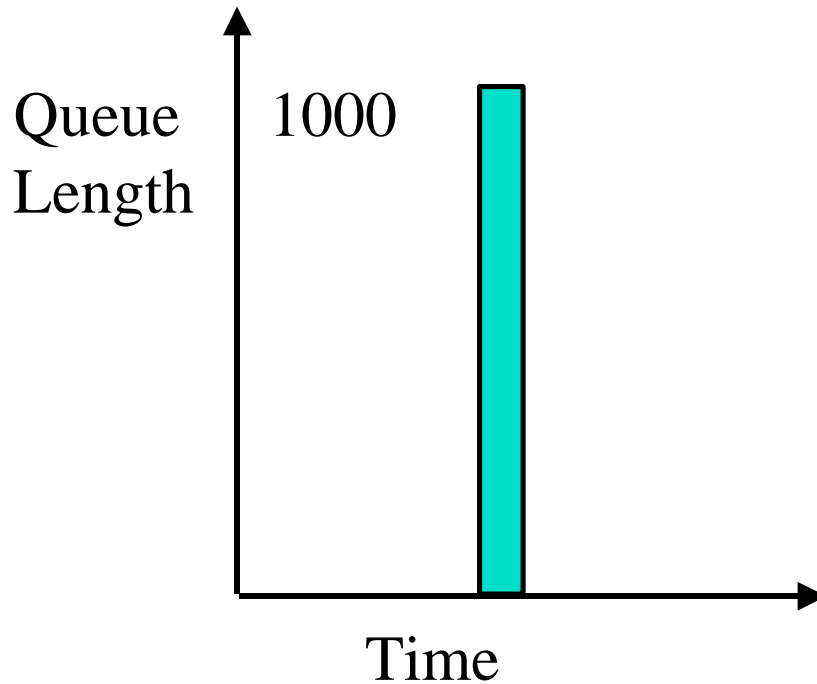
$$= 365,566 \text{ cells/sec} = 155 \text{ Mbps}$$

OSU Congestion Principles

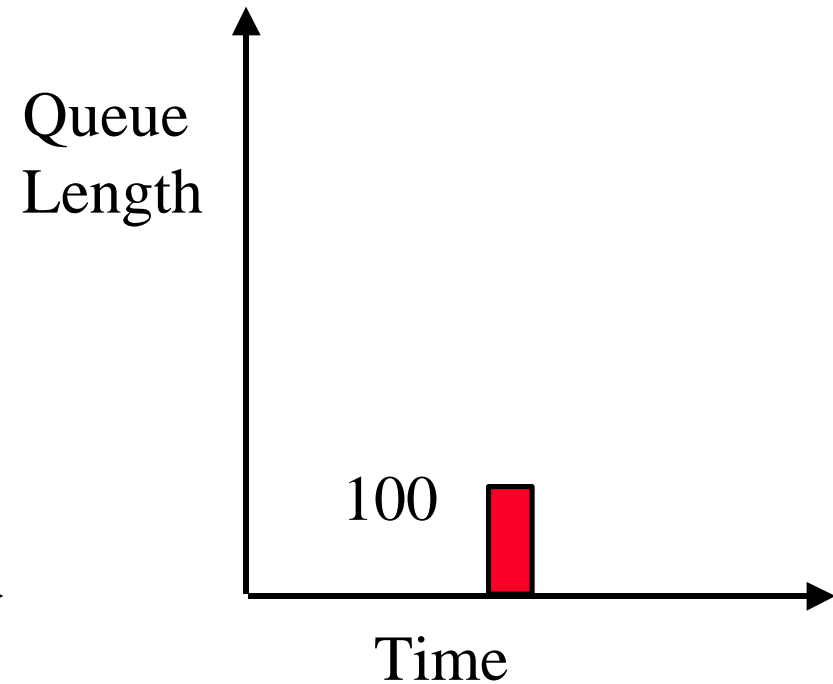
- q Input rate (and not queue length) is the load measure
- q Congestion avoidance (and not congestion control) should be the goal
- q Transient performance (and not the steady state performance) is more important

Which Link is More Overloaded?

Link A

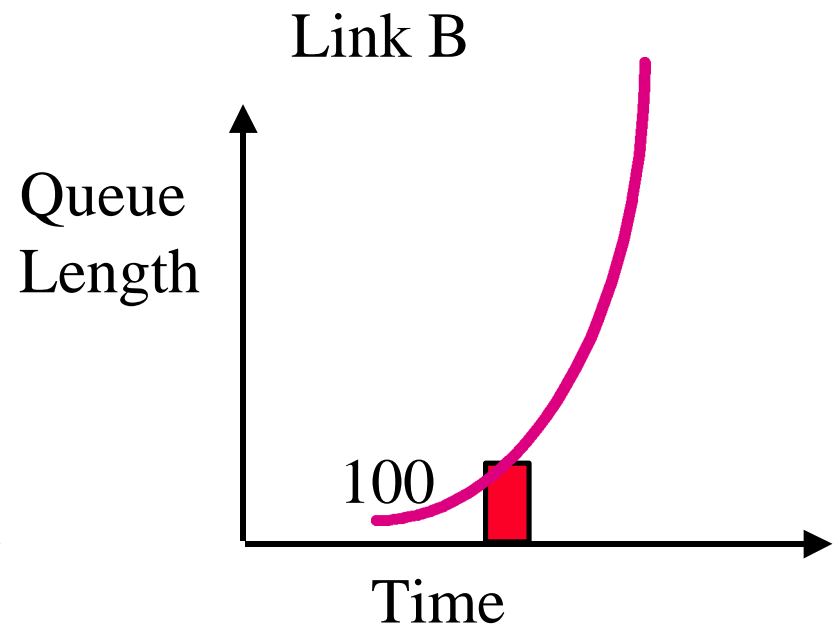
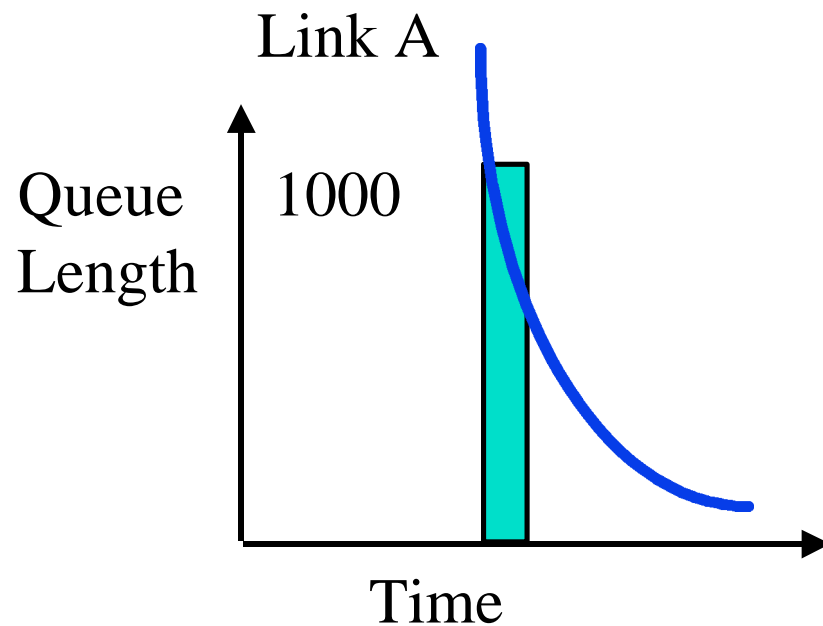


Link B



Answer: It Depends!

- q Link Speed: OC-12 or T1?
- q Control: Rate or Window?
 $Q = \text{Window}$, $dQ/dt = \text{Rate}$
- q For Rate Control: Monitor input rate



Conclusions I

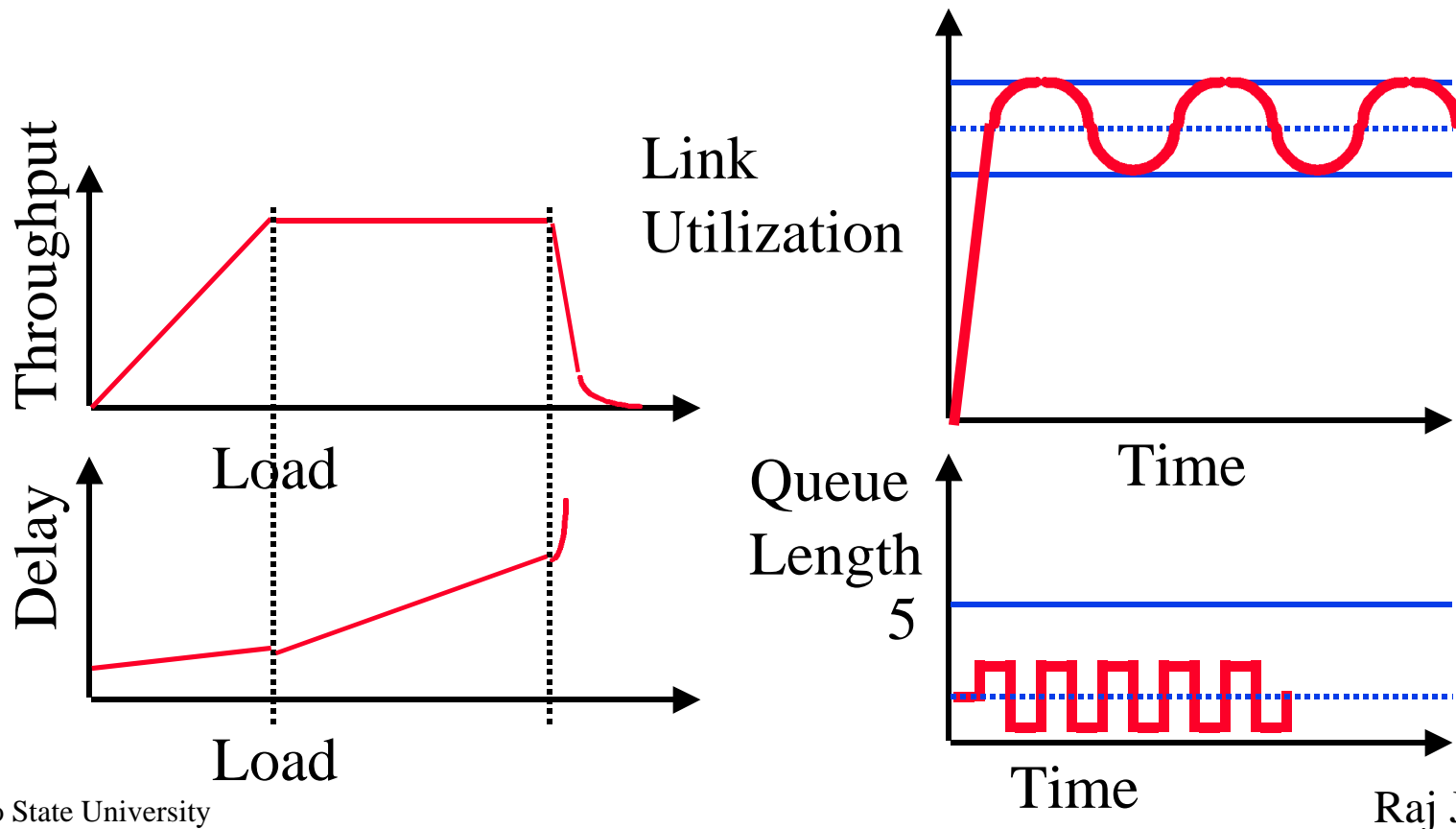
- + Instantaneous queue length is not a good indicator of load for a rate controlled system.

$$Q(t) = Q(t-1) + \text{Input rate}(t) - \text{Service rate}(t)$$

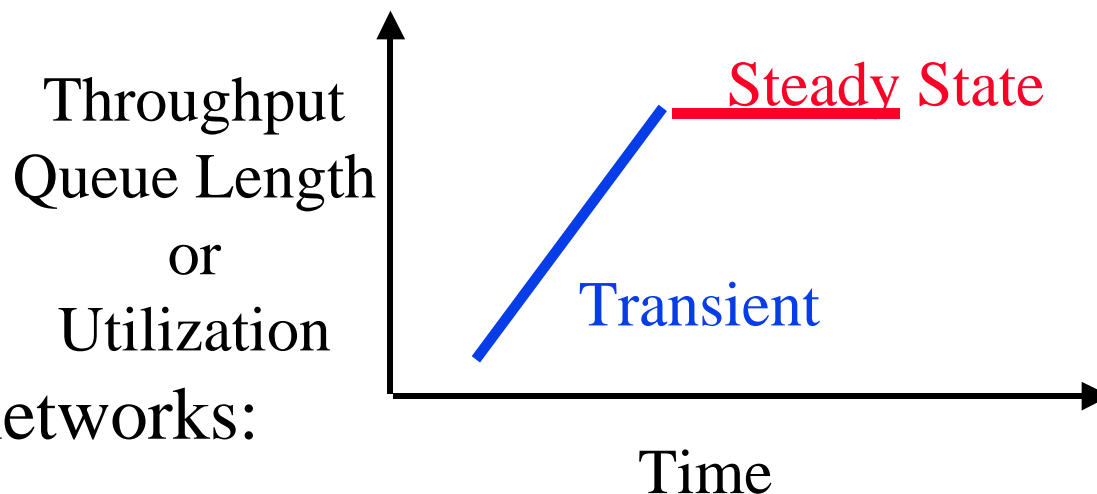
- + Using queue length as the load indicator in a rate controlled system leads to unnecessary oscillations.
- + Input rate monitoring not only correctly tells whether the system is overloaded, it also tells by what factor.
- + Queue = n is not a good goal. Input rate=service rate is.

Congestion Avoidance

- q High throughput, Low delay
- q Small queues
- q Bounded oscillations \Rightarrow Good for Video traffic



Why Worry About Transients?

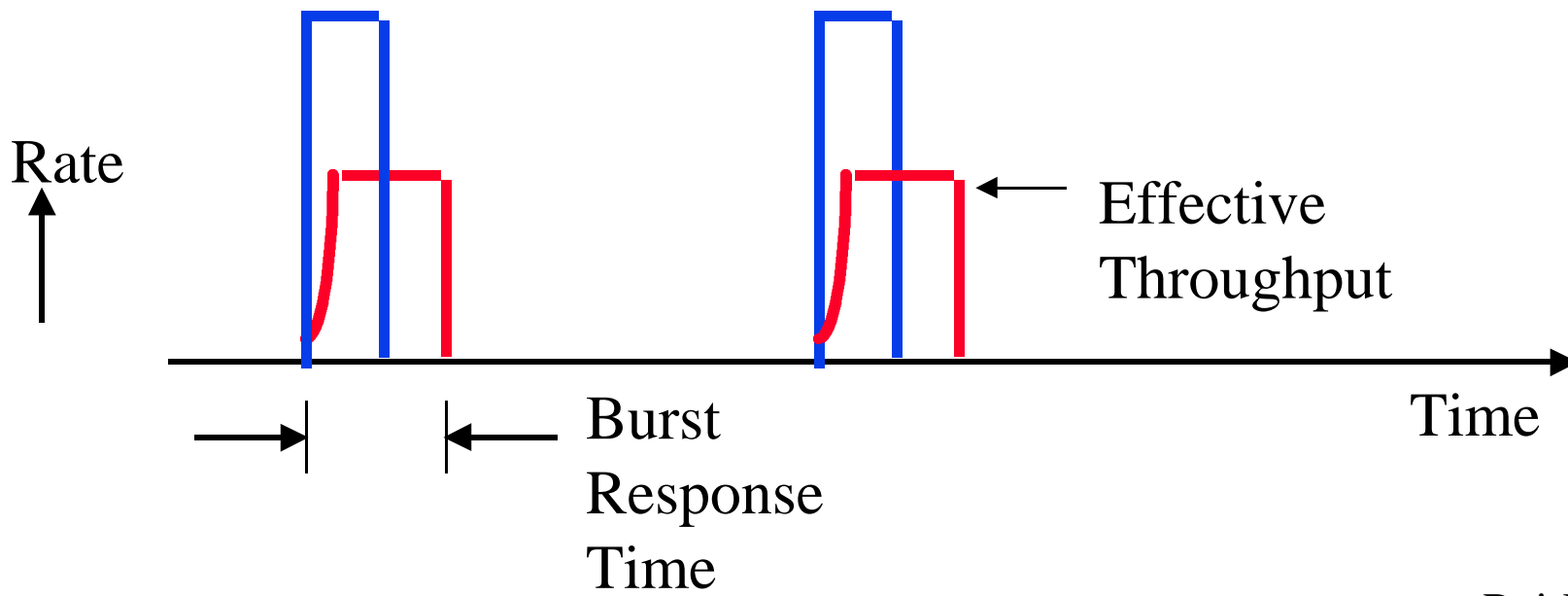
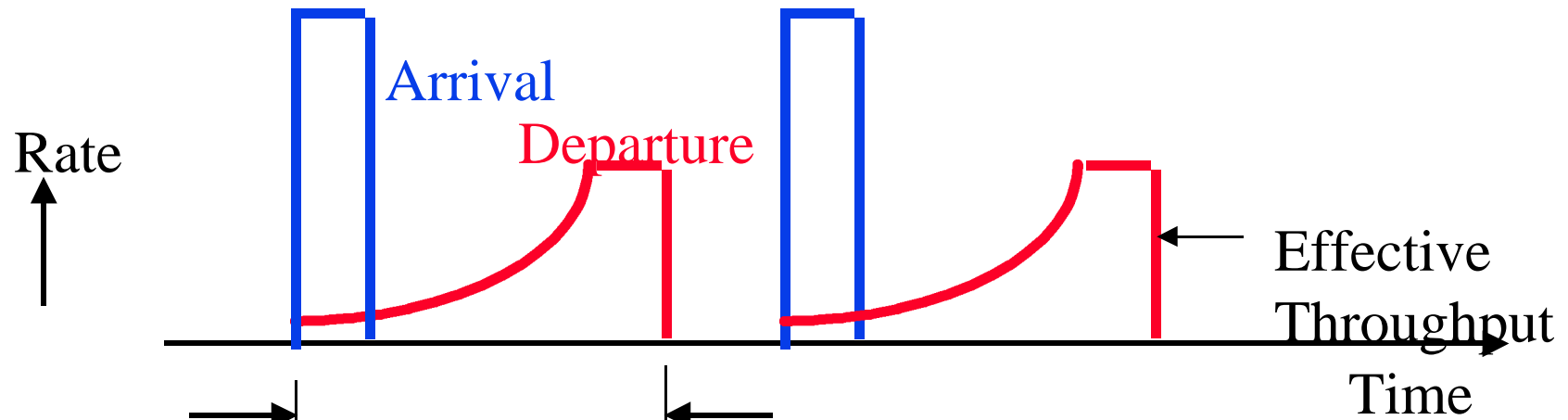


On most networks:

- q There are no infinite sources.
- q Sources come and go
- q VCs may stay but are mostly inactive
- q Traffic is highly bursty

⇒ Networks are operating in the transient region, most of the time.

Burst Performance



Legacy LANs vs ATM

- q Today's LANs have a very fast transient response. Can get to the peak rate within **a few microseconds**
- q On ATM LANs:
Wait for connection setup and then...
Everytime, a burst arrives, take **several milliseconds** to ramp up
- q Q: Given 100 Mbps Switched Ethernet and 155 Mbps ATM at the same price, which one would you buy?

Quiz

T F Please check True/False

1. Congestion is not a problem in high-speed networks
2. User parameter control (UPC) allows a user to set its parameters
3. CDVT measures the cell delay variation caused by the network
4. ABR users do not have to specify CDVT
5. GCRA allows a network to determine conforming and non-conforming cells
6. All non-conforming cells are dropped at the source
7. Credit based scheme requires per-VC queueing
8. EFCI is better than explicit rate for high-speed networks
9. VSVD allows a network operator to use proprietary control scheme inside its networks.

Thank You!

