

1.

9/24

CSE 544T

prop = proportional

Exam I, Oct 8 (Monday)

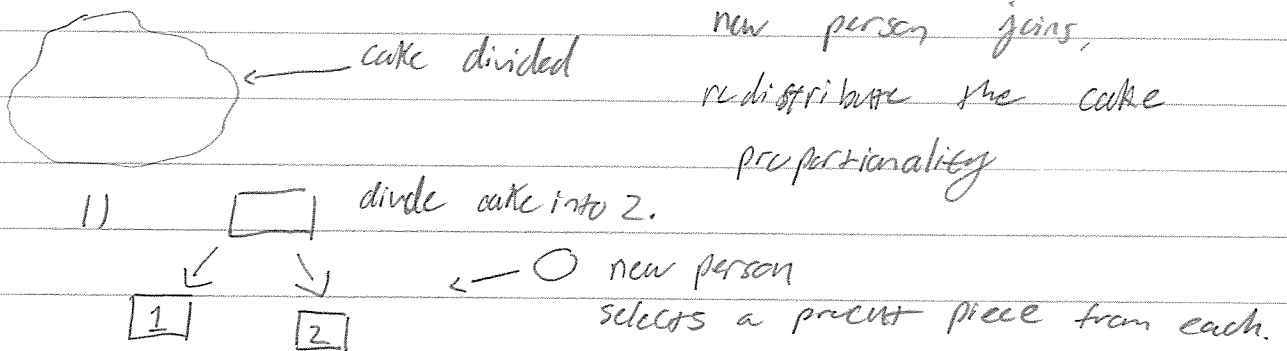
open note, closed internet. Look at * questions

So far:

- * Divide & choose, $n=2$
- Dubins Spainer moving knife*, prop n people
- Austin moving knives; 2 parties decide " $\frac{1}{2}$ "
 - 2 parties decide m " $\frac{1}{m}$ "
 - 2 parties divide cake into m "equal" pieces

- * Branch-Knuster trimming*
 - Fink's lone chooser
 - n parties, dynamic
 - Woodall's strongly prop, $n=2$
 - Steinhaus-Kuhn lone divider prop $n=3$
 - Conway-Guy-Selfridge, $n=3$ envy free

Fink lone chooser:



2.

Fink's lone chooser

start $n=2$ [divide & choose]

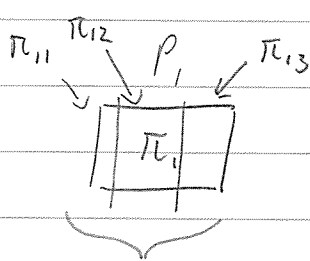
P_1 creates x_1, x_2 where $v(x_1) = v(x_2) = \frac{1}{2}$

Suppose P_2 chooses, takes π_2

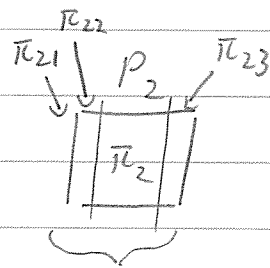
P_1 gets π_1

standard divide & choose

Along comes P_3



equal pieces to
to P_1



equal pieces
to P_2

$$P_1 \quad v(\pi_{11}) = v(\pi_{12}) = v(\pi_{13}) = \left(\frac{1}{2}\right)\left(\frac{1}{3}\right) = \frac{1}{6} \text{ to } P_1$$

$$P_2 \quad v(\pi_{21}) = v(\pi_{22}) = v(\pi_{23}), \text{ each } \geq \frac{1}{6} \text{ to } P_2$$

P_3 gets to pick one piece from π_1^*
and one piece from π_2^*

One piece from P_1 and one from P_2

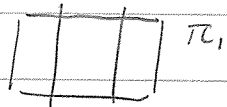
P_3 takes one piece from P_1 so the
value of the remaining cake = $\frac{2}{6} = \frac{1}{3}$, $n=3$
so this is proportional.

3

For P_2 it is the same, value left to $P_2 \geq 2 * \frac{1}{6}$ or $\geq \frac{1}{3}$ $n=3$, proportional.

What about P_3 's value of the cake?

P_1 's piece




value $P_3 = \alpha$

P_2 's piece



value to $P_3 = 1 - \alpha$

if the whole piece of cake has value v , then  to the divider.
3 subpieces

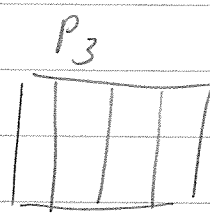
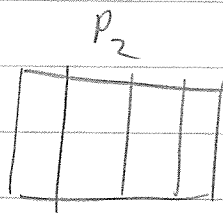
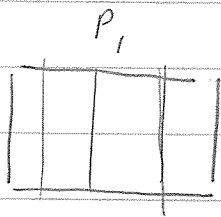
Then not all 3 pieces can have value $< \frac{v}{3}$.

P_3 gets at least $\frac{1}{3} \alpha$ and $\frac{1}{3} (1 - \alpha)$

$$\frac{1}{3} \alpha + \frac{1}{3} (1 - \alpha) = \frac{\alpha + 1 - \alpha}{3} = \frac{1}{3}$$

$\therefore P_3$ picks up cake from P_1 & P_2 with a value of at least $\frac{1}{3}$ from P_3 's perspective

Along comes P_4 divide cake into 4 pieces



This is not envy free.

4

The number of cuts made $O(n^3)$ cuts!

$n =$

$$2 \quad 1$$

$$3 \quad 5 = 4 + 1$$

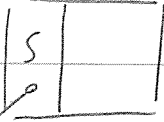
$$4 \quad 14 = 9 + 4 + 1$$

* Woodrall's algorithm Strongly polynomial * $n = 2$
related literature

o Problem of the Nile, 1938 - 1961

o Ham Sandwich theorem, 1942 - 1985

The cakes value must differ to the two parties.

assume cake K 
Diff value \neq $v_1(S) \neq v_2(S)$
2 people \neq $v_1(S) > v_2(S)$

* From the above, show that it is possible
from K to create $\boxed{x_1}$ $\boxed{x_2}$ where x_1 and
 x_2 are necessarily the whole cake.

Such that

$$x_1 \text{ wants } v_1(x_1) > v_1(x_2)$$

$$x_2 \text{ wants } v_2(x_1) < v_2(x_2)$$

Problem

Given x_1, x_2 , & the rest of the cake,

Assign cakes π_1, π_2 such that:

$$V_1(\pi_1) > \frac{1}{2}$$

$$V_2(\pi_2) > \frac{1}{2}$$

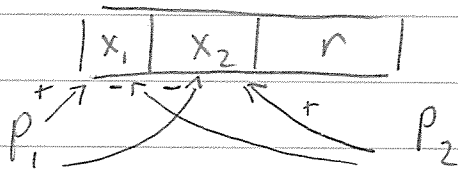
Already have P_1 likes x_1 and P_2 likes x_2 and the rest of the cake, r .

P_1 and P_2 simply play divide & choose on r , creating r_1 & r_2 .

Analysis for person 1 (P_1)

Suppose $V_1(x_1) = \alpha_1, V_1(x_2) = \alpha_2$

$$V_1(r) = 1 - \alpha_1 - \alpha_2$$



By divide and choose P_1 gets r_1 , P_2 gets r_2 and $V(r_1) \geq \frac{1 - \alpha_1 - \alpha_2}{2}$ for P_1 & P_2

Then P_1 takes x_1 and P_2 takes x_2 .

P_1 's
value

$$V(x_1) = \alpha_1$$

$$V(r_1) \geq \frac{1 - \alpha_1 - \alpha_2}{2}$$

$$V(x_1 + r_1) \geq \alpha_1 + \frac{1 - \alpha_1 - \alpha_2}{2} \Rightarrow \frac{2\alpha_1 + 1 - \alpha_1 - \alpha_2}{2}$$

$$V(x_1 + r_1) \geq \frac{1}{2} + \frac{\alpha_1 - \alpha_2}{2}$$

$\therefore V > \frac{1}{2}$ because $\alpha_1 > \alpha_2$

Steinhaus - Kuhn

one divider $n=3$, complicated
example

- ① P_1 creates x_1, x_2, x_3 where $v(x_1) = v(x_2) = v(x_3) = \frac{1}{3}$ from cake
- ② P_2 designates $S_2 \subseteq \{x_1, x_2, x_3\} \mid \forall x \in S_2 \quad v_2(x) \geq \frac{1}{3}$
[lemma $|S_2| \geq 1$]
- ③ P_3 does the same thing to create S_3 using x_1, x_2, x_3

Two possibilities

- a) Lucky $|S_2| > 1$, Assign from smallest to largest $|S_2|, |S_3|$ size.

$$\text{Ex. } S_2 = \{x_1, x_3\} \quad S_3 = \{x_3\}$$

$$\pi_3 = x_3 \quad \pi_2 = x_1 \quad \text{and person}$$

1 gets whatever is left over because they take any piece equally.

- b) Unlucky $|S_2| = |S_3| = 1$

There must be an unacceptable piece (say x_1) such that $v_2(x_1) < \frac{1}{3}$ $v_3(x_1) < \frac{1}{3}$, so give x_1 to P_1 . Then can give x_2 & x_3 and play divide and choose with P_2 & P_3 .

Not envy free. P_1 can envy in case b, (P_2 & P_3 divide & choose)
In case a P_2 can envy P_3 but not P_1 . P_1 cannot envy anyone in a because he values each piece equally.